Roll No.

BT-5/D-21 45170

# FORMAL LANGUAGE AND AUTOMATA THEORY

### Paper-PC-CS-303A

Time Allowed : 3 Hours]

1.

[Maximum Marks : 75

Total Page: 2

Note : Attempt five questions in all, selecting at least one question from each Unit. All questions carry equal marks.

#### UNIT-I

- (a) Prove that the Language A = {0<sup>n</sup> 1<sup>n</sup> | n ≥ 0} is not regular using pumping lemma.
  - (b) Prove that every NFA can be converted to an equivalent DFA that has a single accepting state.
- 2. Give state diagrams of DFAs recognizing the following languages over the alphabet {0, 1}.
  - (a)  $\{W \mid W \text{ contains at least two 0s and at most one 1}\}$ .
  - (b) {W | W starts with 0 and has odd length, or starts with 1 and has even length}.

#### UNIT-II

- (a) Show that the given language {a<sup>i</sup> b<sup>2i</sup> a<sup>i</sup> | i ≥ 0}. is not a CFL using the pumping lemma.
  - (b) Describe the language generated by the CFG with productions  $S \rightarrow ST | \wedge T \rightarrow aS | bT | b$ .
- 4. (a) Let L be the language generated by the CFG with productions  $S \rightarrow aSb|ab|SS$ . Show that no string in L begins with abb.
  - (b) Draw an NFA accepting the language generated by the grammar with productions  $S \rightarrow abA|bB|aba A \rightarrow b|aB|bA B \rightarrow aB|aA$ .

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## UNIT-III

- 5. (a) Give a transition table for PDA that accept the language  $\{a^i b^j | i \le j \le 2i\}$ .
  - (b) Construct a Mealy machine which can generate strings having EVEN and ODD numbers of 1's or 0's.
- 6. (a) Draw a PDA that accept the language :

 $\{0^i \ 1^j \ 2^k \mid i, j, k \ge 0 \text{ and } j = i \text{ or } j = k\}.$ 

(b) Give a transition table for a deterministic PDA that accepts the language  $\{a^i b^{i+j}a^j | i, j \ge 0\}$ .

#### UNIT-IV

- 7. (a) Write down an unrestricted grammar that generate the language  $\{a^n b^n a^n b^n | n \ge 0\}$ .
  - (b) State and explain Cook-Levin theorem.
- (a) Show that the set of languages L over {0, 1} such that neither L nor L' is recursively enumerable is uncountable.
  - (b) Prove that language satisfiable (or the decision problem sat) is NP-complete.